Work Efficient Higher Order Vectorisation

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Data Parallel Haskell (DPH)

• Nested data parallel programming on shared memory multicore.

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- Compiling common classes of programs used to break their asymptotic work complexity.

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- Nested data parallel programming on shared memory multicore.
- Compiling common classes of programs used to break their asymptotic work complexity.
- Not anymore!

User written source code. retrieve :: A (A Char) -> A (A Int) -> A (A Char) retrieve xss iss

= zipWithP mapP (mapP indexP xss) iss

indexP :: A Char -> Int -> Char
mapP indexP xss :: A (Int -> Char)

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nested parallel computation

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builds an array of partially applied functions

indexP :: A Char -> Int -> Char
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Vectorisation

```
User written source code.
retrieve :: A (A Char) -> A (A Int) -> A (A Char)
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        = zipWithP mapP (mapP indexP xss) iss
```

Vectorised data-parallel version.

```
retrieve_v :: A (A Char) -> A (A Int) -> A (A Char)
retrieve_v xss iss
= let ns = takeLengths iss
in unconcat iss
$ index_l (replicates ns xss)
$ concat iss
```

ns = takeLengths **iss** = $[3 \ 1 \ 2 \ 1]$

- ns = takeLengths iss = [3 1 2 1]
- xss1 = replicates ns xss = [[A B] [A B] [A B] [C D E] [F G] [F G] [H]]

ns = takeLengths iss = $[3 \ 1 \ 2 \ 1]$

= [1 0 1 2 1 0 0]

```
ns = takeLengths iss
= [3 \ 1 \ 2 \ 1]
```

xss1 = replicates ns xss = [[A B] [A B] [A B] [C D E] [F G] [F G] [H]] iss2 = concat iss = [1 0 1 2 1 0 0]

 $xss2 = index_l xss1 iss2$ = [B A B E G F H]

ns = takeLengths iss = $[3 \ 1 \ 2 \ 1]$

xss1 = replicates ns xss = [[A B] [A B] [A B] [C D E] [F G] [F G] [H]]

iss2 = concat iss= [1 0 1 2 1 0 0]

xss2 = index_l xss1 iss2
= [B A B E G F H]

res = unconcat **iss xss2** = [[B A B] [E] [G F] [H]]

ns = takeLengths iss = [3 1 2 1]

= [[B A B] [E] [G F] [H]]

xss1 = replicates ns xss = [[A B] [A B] [A B] [C D E] [F G] [F G] [H]] xss1 = replicates [3 1 2 1]
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xss1 = replicates [3 1 2 1]
 [[A B] [C D E] [F G] [H]]
 = [[A B] [A B] [A B] [C D E] [F G] [F G] [H]]

```
retrieve :: A (A Char) -> A (A Int) -> A (A Char)
retrieve xss iss
```

= zipWithP mapP (mapP indexP xss) iss

partial application => sharing

array value [[A B] [A B] [A B] [C D E] [F G] [F G] [H]]

NESL-style array representation



• NESL-style nested array representation cannot represent physical sharing of elements between among sub-arrays.

retrieve [[A B C D E F G H]] [[0 1 2 3 4 5 6 7]] ==> [A B C D E F G H]

- O(size) when using sequential lists (or arrays)
- O(size²) when vectorised. --- BAD!

array value [[A B] [A B] [A B] [C D E] [F G] [F G] [H]]

pointer-based array representation



- Naive pointer based representation has poor locality.
- Hard to distribute array across processors.

Complexity Goal

Suppose the source program is evaluated using pointer-based nested arrays.

The vectorised program should have the same asymptotic work complexity,

but run in parallel,

and be faster than the sequential version.

replicate	::	Int	->	e ->	A e
replicates	::	A Int	-> A	e ->	Ae
concat	::	A (A e)	-> A	е	
unconcat	::	A (A e)	-> A	e ->	A (A e)
pack	::	A Bool	-> A	e ->	Ae
combine2	::	A Bool	-> A	e ->	A e -> A e
index	::	Int	-> A	e ->	е
index_l	::	A (A e)	-> A	Int	-> A e
append	::	Ae	-> A	е	-> A e
append_1	::	A (A e)	-> A	(A e)	-> A (A e)

Pointer Based

replicate	O(length result)
replicates	O(max (len source, len result))
concat	O(max (len source, len result))
unconcat	O(len result)
pack	O(len source)
combine2	O(len result)
index	O(1)
index_l	O(len result)
append	O(len result)
append_1	O(len (concat result))

Old DPH (NESL Style) O(size result) O(max (len source, **size** result)) O(1) O(1) O(max (len source, **size** result)) O(size result) O(1) O(max (len source, **size** result)) O(size result) O(size (concat result))

User written source code.

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Vectorised data-parallel version.

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	Pointer Based	Old DPH (NESL Style)
replicate	O(len result)	O(size result)
replicates	O(max (len source, len result))	O(max (len source, size result))
concat	O(max (len source, len result))	O(1)
unconcat	O(len result)	O(1)
pack	O(len source)	O(max (len source, size result))
combine2	O(len result)	O(size result)
index	O(1)	O(1)
index_l	O(len result)	O(max (len source, size result))
append	O(len result)	O(size result)
append_1	O(len (concat result))	O(size (concat result))

	Pointer Based	New DPH
replicate	O(len result)	O(len result)
replicates	O(max (len source, len result))	O(max (len source, len result))
concat	O(max (len source, len result))	O(max (len source, len result))
unconcat	O(len result)	O(len result)
pack	O(len source)	O(len source)
combine2	O(len result)	O(len result)
index	O(1)	O(1) / O(len result) for nested result
index_l	O(len result)	O(max (len source, len result))
append	O(len result)	O(len result)
append_1	O(len (concat result))	O(len (concat result))

	Pointer Based	New DPH
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pack	O(len source)	O(len source)
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index	O(1)	O(1) / O(len result) for nested result
index_l	O(len result)	O(max (len source, len result))
append	O(len result)	O(len result)
append_l	O(len (concat result))	O(len (concat result))

All implemented using parallel vector operations!

In the paper

- Reference implementation of all operators.
- Other operators (besides replicates) also cause problems.
- Discussion of what complexity operators need to have.
- Invariants needed to maintain complexity.
- How to convert between old and new representations.
- How to avoid using extra descriptor fields when not needed.
- Example arrays and tests

Benchmarks

Sparse Matrix-Vector Multiplication



smvm :: A (A (Int, Double))
 -> A Double -> A Double

smvm matrix vector

= let term (ix, coeff) = coeff * (vector ! ix)
in mapP (\row -> sumP (mapP term row)) matrix



Sparse Matrix-Vector Multiplication (1% non-zero elements)

Barnes-Hut Gravitation Simulation





calcAccels :: Double -> Box -> A MassPoint -> A Accel calcAccels epsilon boundingBox points = mapP (\m -> calcAccel epsilon m tree) points where tree = buildTree boundingBox points Barnes-Hut, 1 step, 1 thread



Things that need to be fixed

- Need array fusion for new representation.
- Integrate with vectorisation avoidance.
- Current parallel implementation is slower than sequential.
- Avoid performing fine grained computations in parallel.

Questions?

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Append Example



replicate	::]	Int	->		e ->	Ae
replicates	::	A	Int	->	A	e ->	Ae
concat	::	A	(A e)	->	A	е	
unconcat	::	A	(A e)	->	A	e ->	A (A e)
pack	::	ΑI	Bool	->	A	e ->	Ae
combine2	::	ΑI	Bool	->	A	e ->	A e -> A e
index	::	Int	t	->	A	e ->	е
index_l	::	A	Int	->	A	(A e)	-> A e
append	::	Αe	9	->	A	е	-> A e
append_1	::	A	(A e)	->	A	(A e)	-> A (A e

)

```
segmap: [0 0 0 1 2 2 3]
source block: [0 0 1 1]
start index: [1 3 0 4]
length: [2 3 2 1]
Blocks 0: [X A B C D E]
1: [F G X X H X X X]
```

```
segmap: [0 0 0 1 2 2 3]
source block: [0 0 1 1]
start index: [1 3 0 4]
length: [2 3 2 1]
Blocks 0: [X A B C D E]
1: [F G X X H X X X]
```

[[K] [] [L M N O]]

```
segmap: [0 1 2]
source block: [0 0 0]
start index: [0 1 1]
    length: [1 0 4]
    Blocks 0: [K L M N 0]
```

```
segmap: [0 0 0 1 2 2 3]
source block: [0 0 1 1]
start index: [1 3 0 4]
length: [2 3 2 1]
Blocks 0: [X A B C D E]
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```

[[K] [] [L M N O]]

segmap: [0 1 2]
source block: [0 0 0]
start index: [0 1 1]
 length: [1 0 4]
 Blocks 0: [K L M N 0]

```
[[A B] [A B] [A B] [C D E] [F G] [F G] [H] [K] [] [L M N O]]
segmap: [0 0 0 1 2 2 3 4 5 6]
source block: [0 0 1 1 2 2 2]
start index: [1 3 0 4 0 1 1]
length: [2 3 2 1 1 0 4]
Blocks 0: [X A B C D E]
1: [F G X X H X X X]
2: [K L M N O]
```

segmap: [0 0 0 1 5 5 6]
source block: [0 0 1 1 0 1 1 1 1]
start index: [1 3 2 2 0 5 7 0 4]
length: [2 3 2 1 2 1 2 1]
Blocks 0: [X A B C D E]
1: [F G X X H X X]

[[K] [] [L M N O]]

segmap: [0 1 2]
source block: [0 0 0]
start index: [0 1 1]
 length: [1 0 4]
 Blocks 0: [K L M N 0]

```
[[A B] [A B] [A B] [C D E] [F G] [F G] [H] [K] [] [L M N O]]
segmap: [0 0 0 1 5 5 6 7 8 9]
source block: [0 0 1 1 0 1 1 1 1 2 2 2]
start index: [1 3 2 2 0 5 7 0 4 0 1 1]
length: [2 3 2 2 1 2 1 2 1 2 1 4 0 4]
Blocks 0: [X A B C D E]
1: [F G X X H X X X]
2: [K L M N O]
```

segmap: [0 0 0 1 5 5 6]
source block: [0 0 1 1 0 1 1 1 1]
start index: [1 3 2 2 0 5 7 0 4]
length: [2 3 2 1 2 1 2 1]
Blocks 0: [X A B C D E]
1: [F G X X H X X]

[[K] [] [L M N O]]

segmap: [0 1 2]
source block: [0 0 0]
start index: [0 1 1]
 length: [1 0 4]
 Blocks 0: [K L M N 0]

Invariant

All physical segment descriptors must be reachable from the virtual segment map

```
[[A B] [A B] [A B] [C D E] [F G] [F G] [H] [K] [] [L M N O]]
segmap: [0 0 0 1 5 5 6 7 8 9]
source block: [0 0 1 1 0 1 1 1 1 2 2 2 ]
start index: [1 3 2 2 0 5 7 0 4 0 1 1]
length: [2 3 2 2 1 2 1 2 1 2 1 0 4]
Blocks 0: [X A B C D E]
1: [F G X X H X X X]
2: [K L M N O]
```

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Tree Lookup Benchmark

```
treeLookup :: A Int -> A Int -> A Int
treeLookup table indices
  | lengthP indices == 1
  = [:table !: (indices !: 0):]
  | otherwise
  = let half = lengthP indices `div` 2
     s1 = sliceP 0 half indices
     s2 = sliceP half half indices
     in concatP (mapP (treeLookup table) [: s1, s2 :])
```

Tree Lookup



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Index Space Overflow

```
retsum :: A (A Int) -> A (A Int) -> A (A Int)
retsum xss iss
  = zipWithP mapP
      (\xs i. indexP xs i + sumP xs) xss) iss
```

```
retsum [[1 2] [4 5 6] [8]] (xss)
[[1 0 1] [1 2] [0]] (iss)
==> [[5 4 5] [20 21] [16]]
```

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Rewriting to Simplified Segds



Backing off to simpler segment descriptors

- sum_vs :: VSegd -> PDatas Int -> PData Int
- sum s :: Segd -> PData Int -> PData Int

- RULE "sum_vs/promote" forall segd arr
 - . sum_vs (promoteSSegd (promoteSegd segd))
 (singletondPR arr)
 - = sum_s segd arr

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Space Complexity

```
furthest :: PA (Float, Float) -> Float
furthest ps = maxP (mapP (\p. maxP (mapP (dist p) ps)) ps)
```